



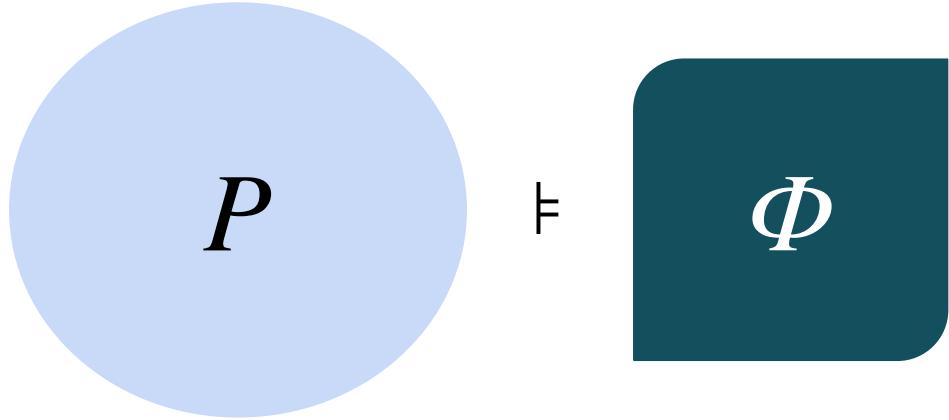
Computing Precise Control Interface Specifications

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The Unknown

How do we verify programs with unknown code?

A Verification Myth



Program

Spec

Source code is incomplete!

Libraries

Modules

System
Calls

$$F = \{F_1, \dots, F_n\}$$

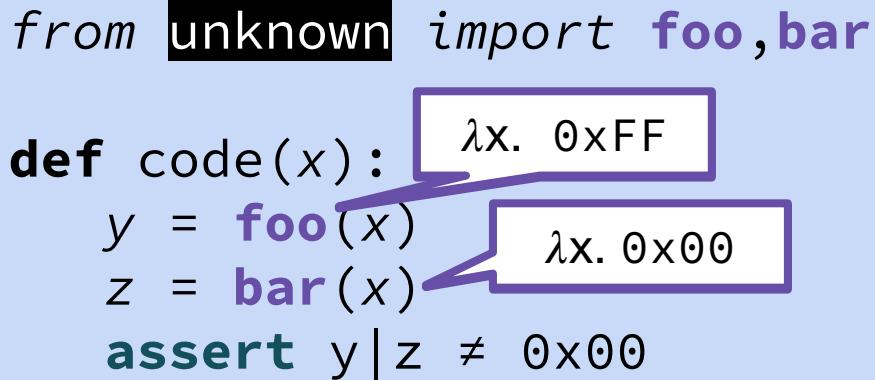
$P[F]$

Program

\models Φ

Spec

```
from unknown import foo,bar  
  
def code(x):  
    y = foo(x)  
    z = bar(x)  
    assert y|z != 0x00
```



The diagram illustrates the flow of control from the function definitions to the calls in the code. Two purple speech bubbles contain lambda expressions: $\lambda x. 0xFF$ above the call to `foo(x)` and $\lambda x. 0x00$ above the call to `bar(x)`. Arrows point from these bubbles to their respective function calls in the code.

Which implementations?

```
import unknown as foo,bar  
  
def code(x):  
    y = foo(x)  
    z = bar(x)  
    assert y|z != 0x00
```

Which implementations satisfy the spec?

Independent Specs
[POPL '16, SIGCOMM '20]

$\varphi(\text{foo}) \wedge \psi(\text{bar})$

Necessary [SIGCOMM '20]

true

Not Safe!!!

Eliminates no “good runs”
[VMCAI '13]

Safe

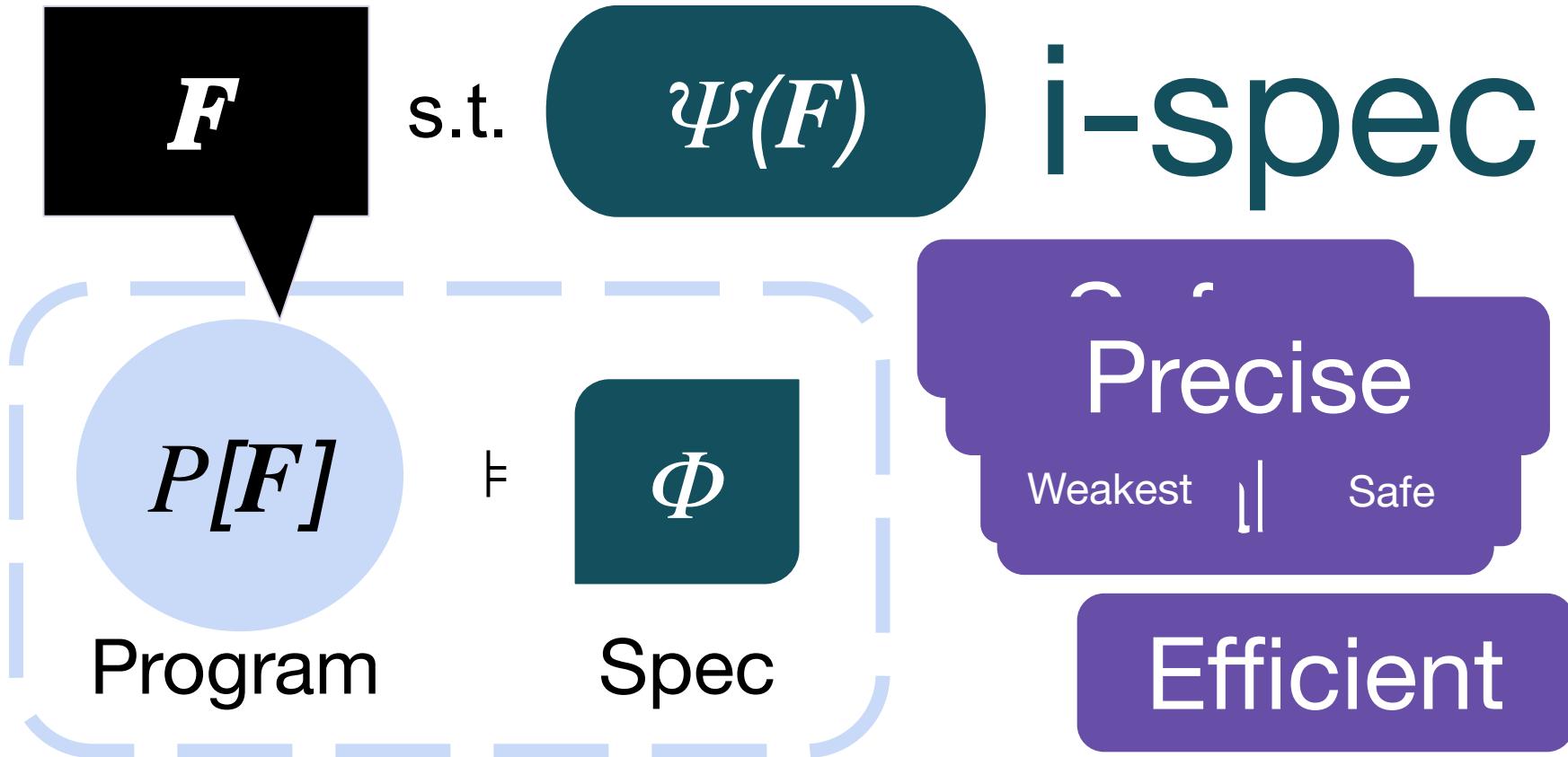
Unsafe!

Overly
Restrictive!

Permissive

[POPL '16]

[SIGCOMM '20]



```
import unknown as foo,bar  
  
def code(x):  
    y = foo(x)  
    z = bar(x)  
    assert y|z != 0x00
```

Goal:
Compute
precise i-specs

foo(x) | bar(x) ≠ 0x00

Safe

Weakest

Efficient

How to use computed i-specs?

Documentation

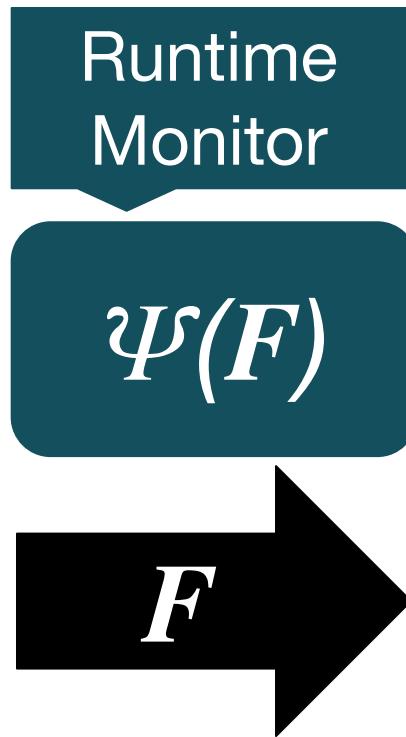


Manual
Verification



Automated
Verification

Control
Plane
(e.g. BGP)



Data Plane
Pipeline
(e.g. P4)

Computer Network

Capisce

computes control interface specs
(ci-specs)

Precise

Safe

Efficient

Weakest

Step 1: Model Pipeline Programs

GCL(\mathbf{F})

?

c ::= assume φ
| x := e
| c ; c
| c [] c

$\mathbf{F} : 2^w \rightarrow 2^l$

e ::= ... | $\mathbf{F}(e)$

$\varphi \in \text{QFBV}$

Control Flow in GCL(\mathbf{F})

```
if  $\varphi$  c1 c2
```

```
assume  $\varphi$ ; c1
      []
assume  $\neg\varphi$ ; c2
```

Step 2: Symbolic Compilation

GCL(F)

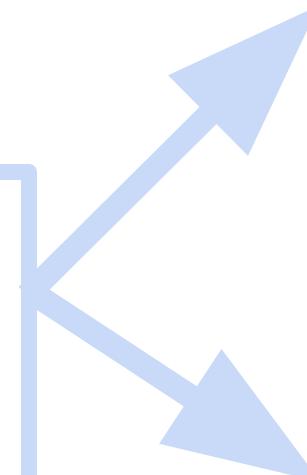
Lifting

assume
 $\vartheta(F, \text{cfg})$;

$c[F, \text{pkt}]$

$c'[\text{cfg}, \text{pkt}]$

GCL

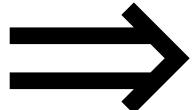


assume $\vartheta(\mathbf{F}, \mathbf{cfg})$;

$c'[\text{pkt}, \mathbf{cfg}]$

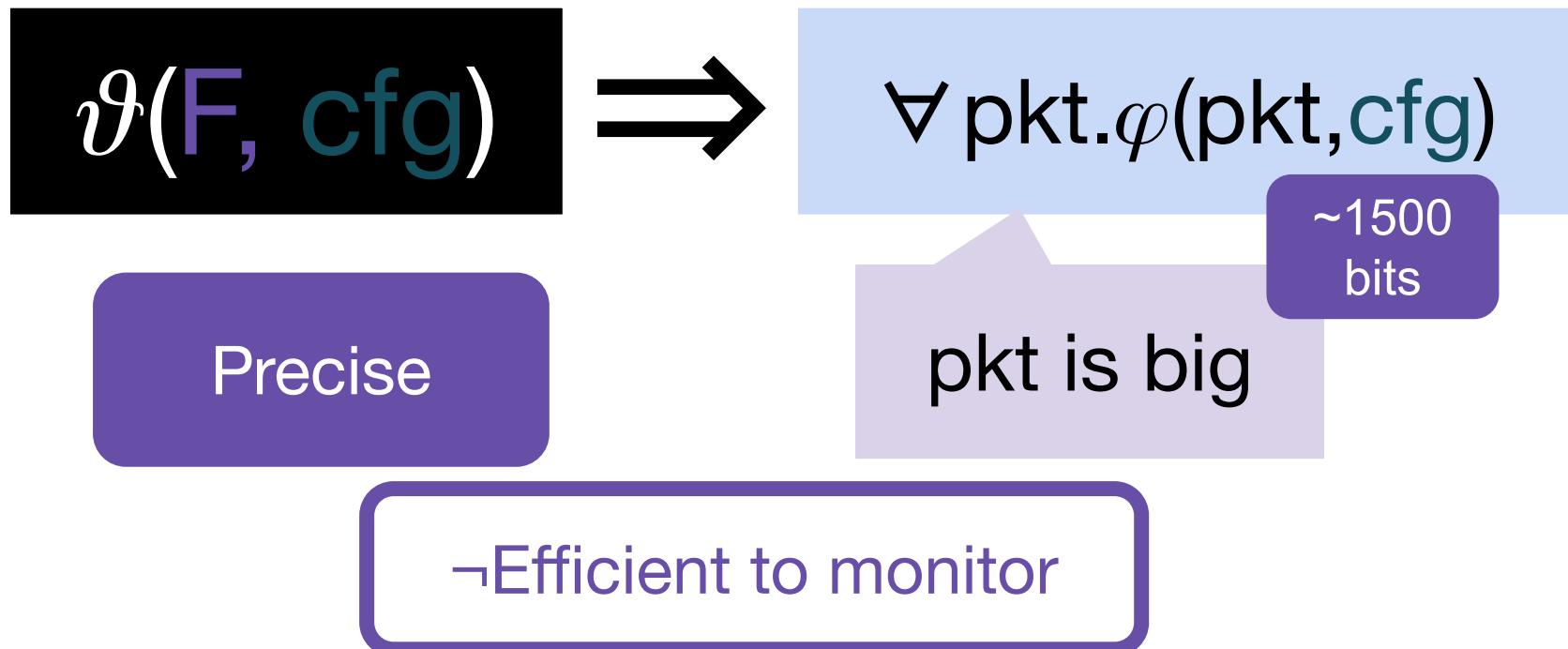
Symbolic
Compilation

$\vartheta(\mathbf{F}, \mathbf{cfg})$



$\forall \text{pkt}. \varphi(\text{pkt}, \mathbf{cfg})$

ci-spec



Step 3: Quantifier Elimination

$$\vartheta(\mathsf{F}, \text{cfg}) \implies \forall \text{pkt}. \varphi(\text{pkt}, \text{cfg})$$

QE

$$\vartheta(\mathsf{F}, \text{cfg}) \implies \psi(\text{cfg})$$

ci-spec

$$\vartheta(F, y) \implies \psi(y)$$

Theorem. *Precise*
 \Rightarrow *safe*
 \Rightarrow *weakest*

Theorem. *Efficiently monitorable*

Efficiently Control-Monitorable Sentences

Theorem. *Terminates*

... have polynomial expression complexity

Quantifier
Elimination
is *Intractable*

$\forall \text{pkt.} \varphi(\text{pkt}, \text{cfg})$

QE

$\psi(\text{cfg})$

new candidate ci-spec

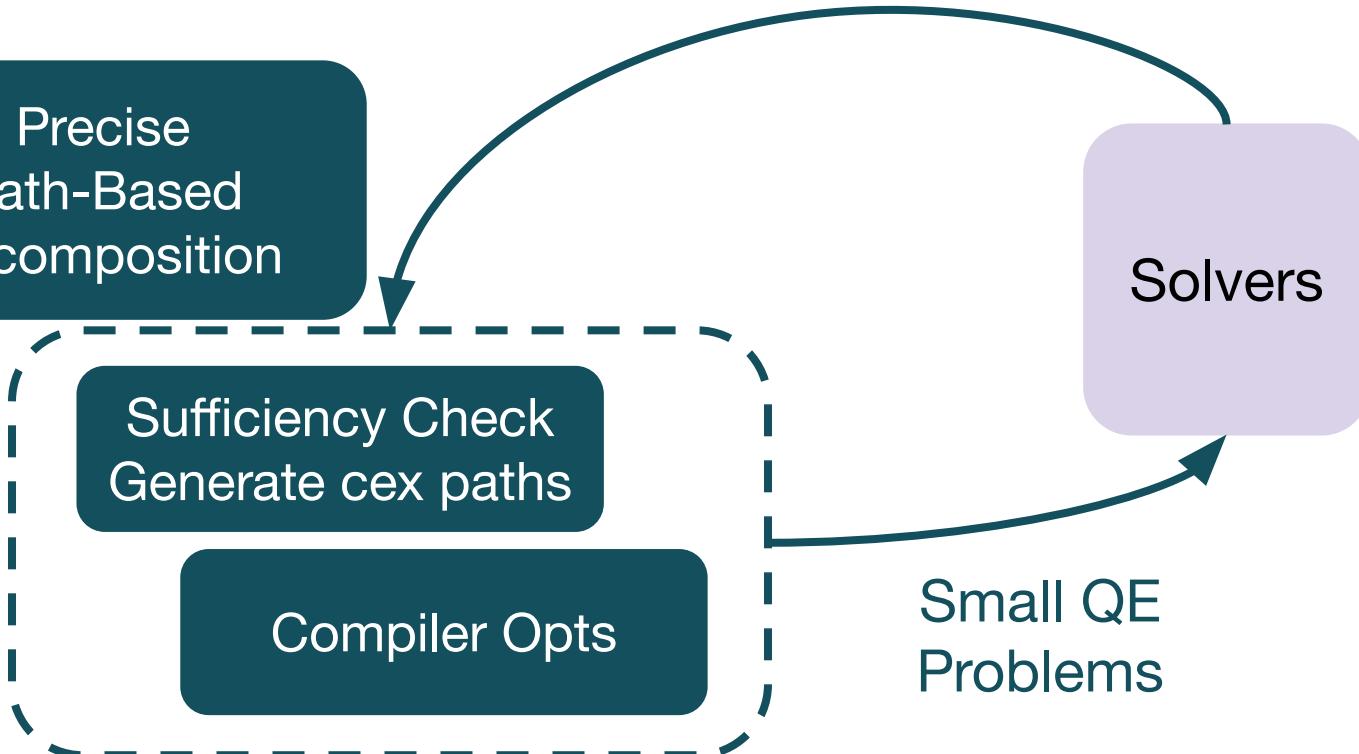
Precise
Path-Based
Decomposition

Sufficiency Check
Generate cex paths

Compiler Opt

Solvers

Small QE
Problems



Evaluation

Survey of Industrial and Academic P4 Programs

Ensure Invalid Data is not Read

Program	Program Paths	Result	Time (s)	Explored Paths	Spec AST Size	Explored Ratio
ABURD PROGRAMS						
ts-switching	21	⊥	0.160	2	1	0.095
mc-nat	39	⊥	0.089	1	1	0.026
FIXES TO ABSURD PROGRAMS						
ts-switching-fixed	21	⊤	0.030	0	1	0.0
mc-nat-fixed	39	⊤	0.027	0	1	0.0
TRIVIAL PROGRAMS						
resubmit	9	⊤	0.028	0	1	0.0
netpaxos-acceptor	0.116	⊤	30.0	0	1	0.0
ecmp	102	⊤	0.030	0	1	0.0
hula	3629	⊤	0.068	0	1	0.0
ndp-router	3843	⊤	2.9	0	1	0.0
NON-TRIVIAL PROGRAMS						
arp	95	∅	5.0	0.016	349	0.17
heavy-hitter-2	267	∅	0.29	3	26	0.011
heavy-hitter-1	327	∅	0.60	7	90	0.021
flowlet	649	∅	1.8	9	127	0.014
simple_nat	66531	∅	5.2	54	1421	0.00081
07-multiprotocol	54459	∅	16	143	3138	0.0026
netchain	26726780	∅	2.9×10^3	264	11658	9.9×10^{-6}
linearroad	54477696	∅	timeout			
fabric	133365047559893	∅	timeout			
SPEC SENSITIVITY PROGRAMS						
heavy-hitter-1-fixed	327	∅	0.63	7	107	0.021
linearroad-fixed	54477696	∅	5.9×10^4	3236	179885	5.9×10^{-5}
fabric-fixed	133365047559893	∅	1.2×10^3	653	41140	4.9×10^{-12}

Program Survey

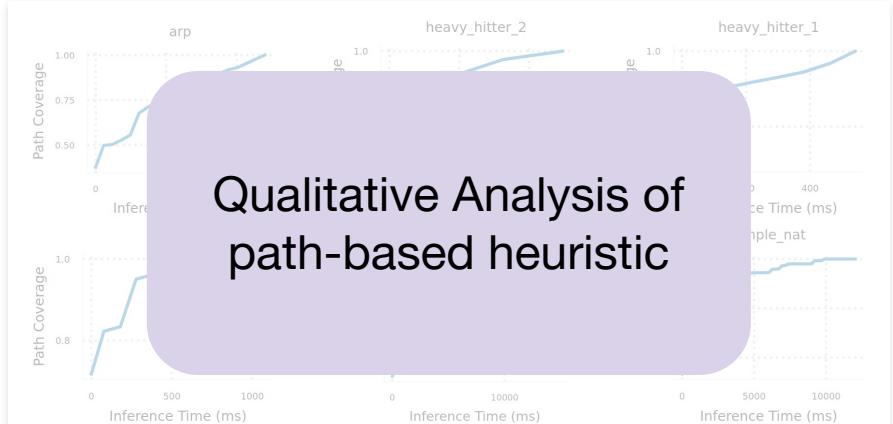
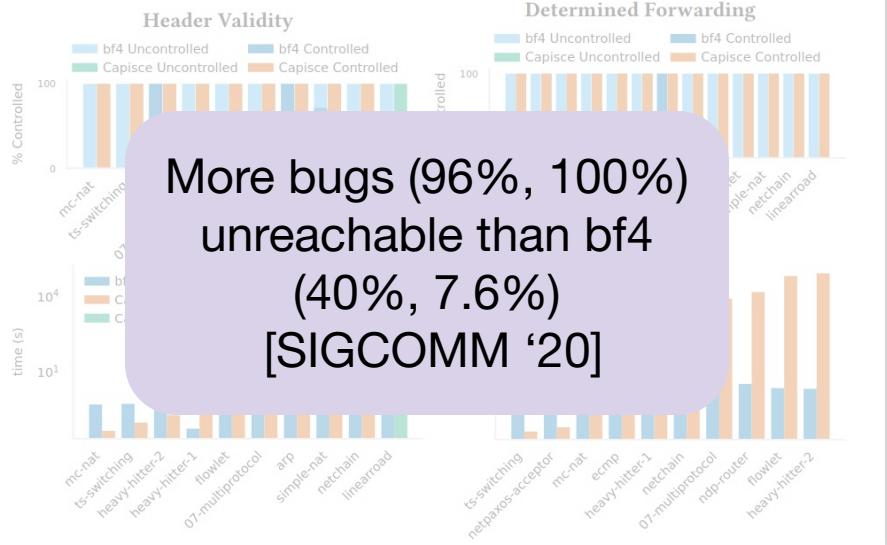
Ensure Invalid
Data is not
Read

Program	Program Paths	Result	Time (s)	Explored Paths	Spec AST Size	Explored Ratio
ABSURD PROGRAMS						
ts-switching	21	⊥	0.160	2	1	0.095
mc-nat	39	⊥	0.089	1	1	0.026
FIXES TO ABSURD PROGRAMS						
ts-switching-fixed	21	⊤	0.030	0	1	0.0
mc-nat-fixed	39	⊤	0.027	0	1	0.0
TRIVIAL PROGRAMS						
resubmit	9	⊤	0.028	0	1	0.0
netpaxos-acceptor	0.116	⊤	30.0	0	1	0.0
ecmp	102	⊤	0.030	0	1	0.0
hula	3629	⊤	0.068	0	1	0.0
ndp-router	3843	⊤	2.9	0	1	0.0
NONTRIVIAL PROGRAMS						
arp	95	φ	5.0	0.016	349	0.17
heavy-hitter-2	267	φ	0.29	3	26	0.011
heavy-hitter-1	327	φ	0.60	7	90	0.021
flowlet	649	φ	1.8	9	127	0.014
simple_nat	66531	φ	5.2	54	1421	0.00081
07-multiprotocol	54459	φ	16	143	3138	0.0026
				264	11658	9.9×10^{-6}
linearroad	54477696		timeout			
fabric	133365047559893		timeout			
SPEC SMELL PROGRAM FIXES						
linearroad-fixed	54477696	φ	5.9×10^4	3236	179885	5.9×10^{-5}
fabric-fixed	133365047559893	φ	1.2×10^3	653	41140	4.9×10^{-12}

Program Survey

Defined
Forwarding

Program	Program Paths	Result	Time (s)	Explored Paths	ci-spec Size	Explored Ratio
		Abs	ARD PROGRAMS			
ecmp	102	⊥	0.320	4	1	0.039
fabric	133365047559893	⊥	7.3	5	1	3.7×10^{-14}
netchain	26726780	⊥	27	7	1	2.6×10^{-7}
		Triv	TRIVIAL PROGRAMS			
arp	95	⊤	0.027	0	1	0.0
linearroad	54477696	⊤	0.054	0	1	0.0
simple-nat	5548	⊤	0.034	0	1	0.0
		Nont	NONTRIVIAL PROGRAMS			
resubmit	9	φ	0.016	2	17	0.22
ts-switching	21	φ	0.10	1	4	0.048
mc-nat	39	φ	0.27	3	21	0.077
netpaxos-acceptor	116	φ	0.12	1	4	0.0086
heavy-hitter-2	267	φ	88	15	233	0.056
heavy-hitter-1	327	φ	0.10	11	187	0.034
flowlet	649	φ	79	15	490	0.023
hula	3629	φ	0.39	1	9	0.00028
ndp-router	3843	φ	40	36	824	0.0094
07-multiprotocol	54459	φ	30	232	5034	0.0043
		Spec	SPECIMENTS & FINES			
ecmp-fixed	102	φ	0.28	3	34	0.029
mc-nat-fixed	27	⊤	0.029	0	1	0.0



Capisce

$\psi(F)$

$p[F]$

ϕ

F

Guarded Pipeline Language (GPL)

p ::= $\mathbf{t}(e)$ *table*
| c $\in GCL$



e ::= ... φ ::= ...

$\mathbf{t} : 2^w \rightarrow \{a_1, \dots, a_n\}$

a ::= $\lambda x. c$ $c \in GCL$

Declaration

```
fwd : 232 → {  
  (λp:29. port := p),  
  (λ_:20. drop := 1)  
}
```

Program

```
fwd(ipv4_dst)
```

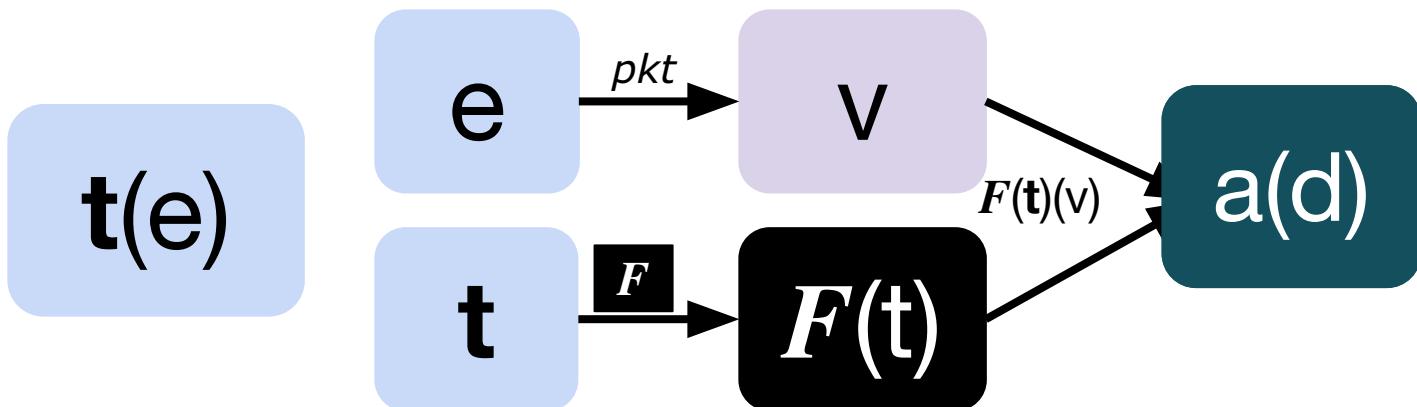
Unknown
Implementation

$F : \text{Tbl} \rightarrow \text{BV} \rightarrow \text{Act} \times \text{BV}$

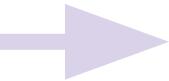
Variable Store

$pkt : \text{Var} \rightarrow \text{BV}$

Table
Semantics



GPL



GCL(F)

```
fwd :  $2^{32} \rightarrow \{$ 
   $(\lambda p:2^9. \text{ port} := p),$ 
   $(\lambda _:2^0. \text{ drop} := 1)$ 
}
```

```
fwd(ipv4_dst)
```



```
Fwd :  $2^{32} \rightarrow 2^1 \times 2^9$ 
i, p := Fwd(ipv4_dst);
if i = 0 {
  port := p
} else {
  drop := 1
}
```

P4

// table declarations
forward : bitvec<32> -> {
 (lambda dmac : bitvec<48>. hdr.ethernet.dstAddr := dmac),
 (lambda _ : bitvec<0>. standard_metadata.egress_spec := 511),
 (lambda _ : bitvec<0>. assume true)
}

send_frame : bitvec<9> -> {
 (lambda smac : bitvec<48>. hdr.ethernet.srcAddr := smac),
 (lambda _ : bitvec<0>. standard_metadata.egress_spec := 511)
}

ecmp_group : bitvec<48> -> {
 (lambda _ : bitvec<0>. standard_metadata.egress_spec 511),
 (lambda (nhop : bitvec<32>). (port : bitvec<9>).
 meta.routing_metadata.nhop_ipv4 := nhop_ipv4,
 standard_metadata.egress_spec := port;
 hdr.ipv4.ttl := ipv4.ttl - 1);
 (lambda _ : bitvec<0>. assume true)
}

zombie.parse_result := 0;
hdr.ethernet.isValid ;
if (var hdr.ethernet.ethertype == 0x0800){
 hdr.ipv4.isValid := 1;
 if (hdr.ipv4.protocol == 6){
 hdr.tcp.isValid := 1;
 zombie.parse_result := 1;
 } else {
 zombie.parse_result := 1;
 }
} else {
 zombie.parse_result := 1;
};
if (hdr.ipv4.isValid = 1)){
 if (hdr.ipv4.ttl > 0){
 ecmp_group(hdr.ipv4.dst);
 forward(meta.routing_metadata.nhop_ipv4));
 }
}
if (standard_metadata.egress_spec != 511){
 standard_metadata.egress_port := standard_metadata.egress_port;
 send_frame(standard_metadata.egress_port)
}